

Networked Applications: Sockets

Fall 1389

Acknowledgments: Lecture slides are from Computer networks course thought by Jennifer Rexford at Princeton University. This presentation was edited for CE443 by Sadeqh Dorri <dorri@ce.sharif.edu> and Behnam Momeni <b_momeni@ce.sharif.edu>.

Goals of Today's Lecture

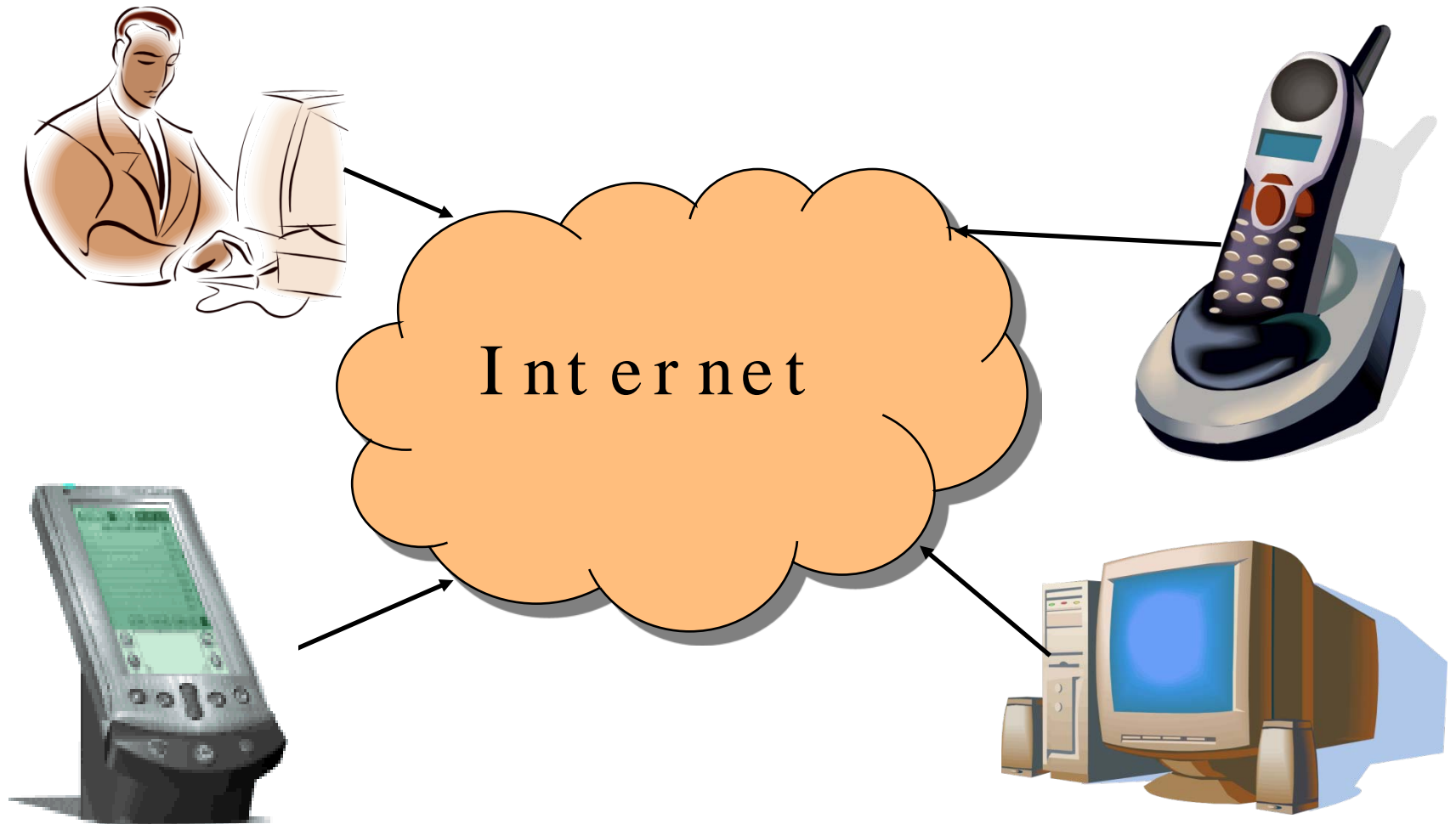


- Client-server paradigm
 - End systems
 - Clients and servers
- Sockets
 - Socket abstraction
 - Socket programming in UNIX
- HyperText Transfer Protocol (HTTP)
 - URL, HTML, and HTTP
 - Clients, and servers
 - Example transactions using sockets



Client-Server Paradigm

End System: Computer on the 'Net



Also known as a “host”...

Clients and Servers



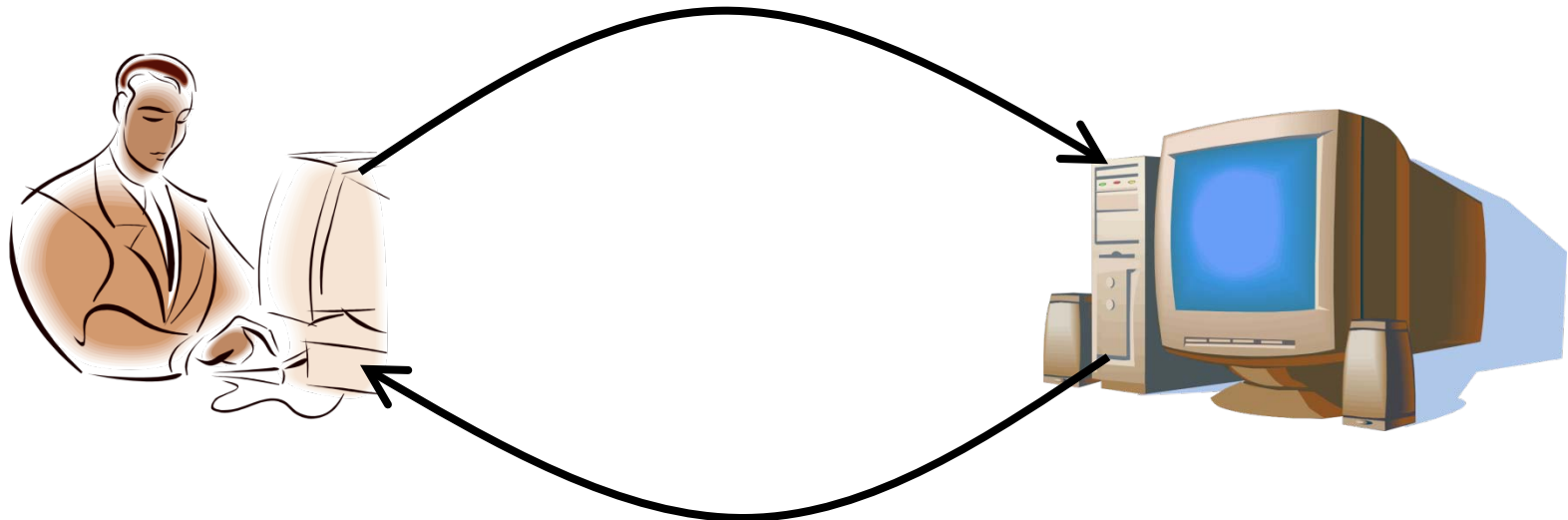
- Client program

- Running on end host
- Requests service
- E.g., Web browser

- Server program

- Running on end host
- Provides service
- E.g., Web server

GET / i n d e x . h t m l

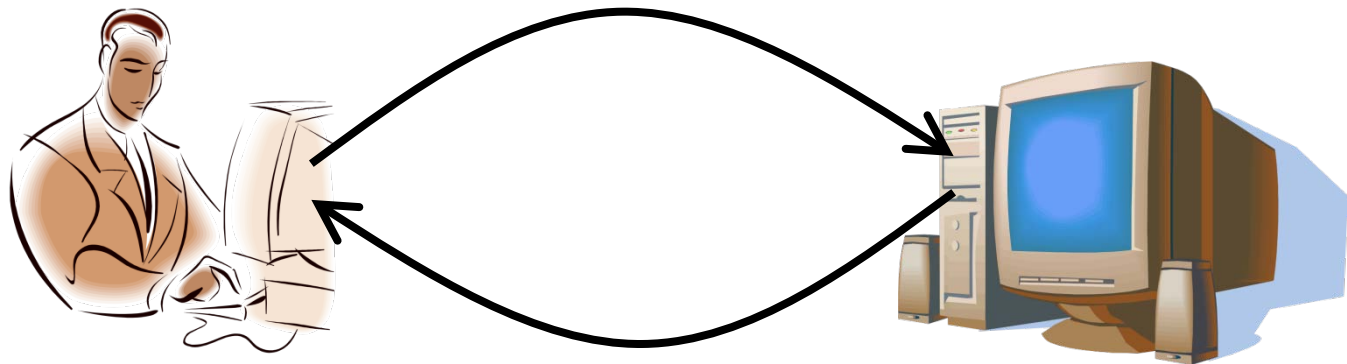


“Site under construction”

Client-Server Communication



- Client “sometimes on”
 - Initiates a request to the server when interested
 - E.g., Web browser on your laptop or cell phone
 - Doesn’t communicate directly with other clients
 - Needs to know the server’s address
- Server is “always on”
 - Services requests from many client hosts
 - E.g., Web server for the www.cnn.com Web site
 - Doesn’t initiate contact with the clients
 - Needs a fixed, well-known address



Peer-to-Peer Paradigm



- No always-on server at the center of it all
 - Hosts can come and go, and change addresses
 - Hosts may have a different address each time
- Example: peer-to-peer file sharing
 - Any host can request files, send files, query to find a file's location, respond to queries, ...
 - Scalability by harnessing millions of peers
 - Each peer acting as both a client and server

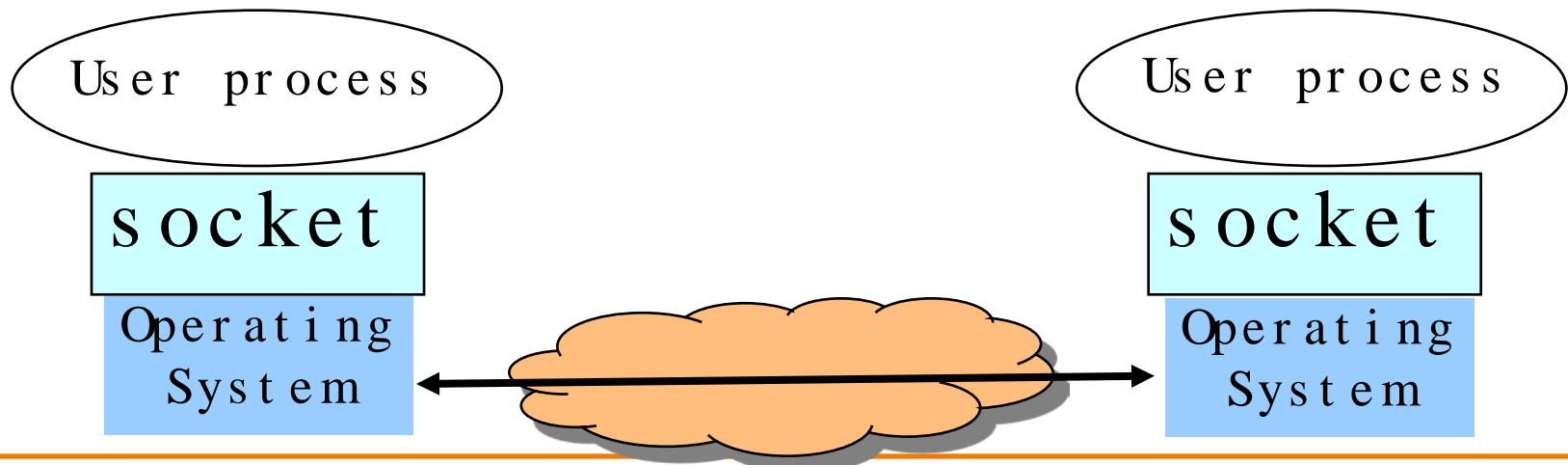


Sockets

Socket: End Point of Net. Comm.'s



- Socket as an Application Programming Interface
 - Supports the creation of network applications
- Two ends communicate through a “socket”
 - Sending messages from one process to another
 - The transportation details are transparent to the programmer



Delivering the Data: Division of Labor



- Application

- Read data from and write data to the socket
- Interpret the data (e.g., render a Web page)

- Operating system

- Deliver data to the destination socket
- Based on the destination port number



- Network

- Deliver data packet to the destination host
- Based on the destination IP address

Identifying the Receiving Process



- Sending process must identify the receiver
 - The receiving end host machine
 - The specific socket in a process on that machine
- Receiving host
 - Destination address that uniquely identifies the host
 - An IPv4 address is a 32-bit quantity
- Receiving socket
 - Host may be running many different processes
 - Destination port that uniquely identifies the socket
 - A port number is a 16-bit quantity

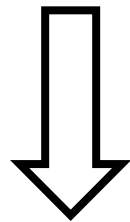
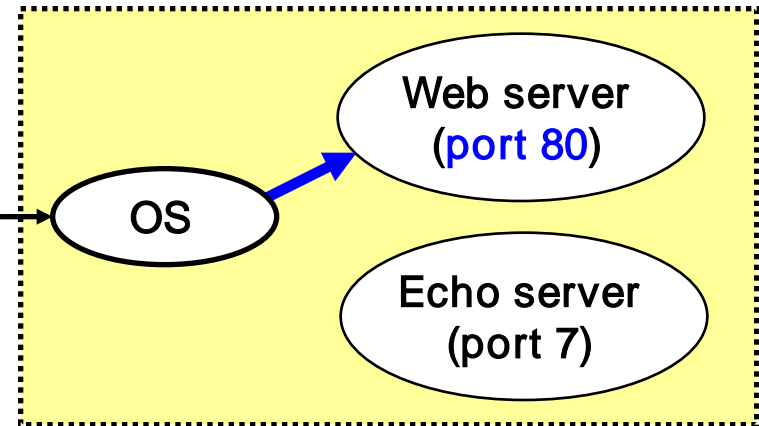
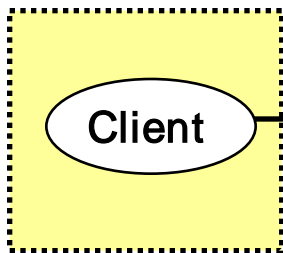
Identifying the Receiving Process



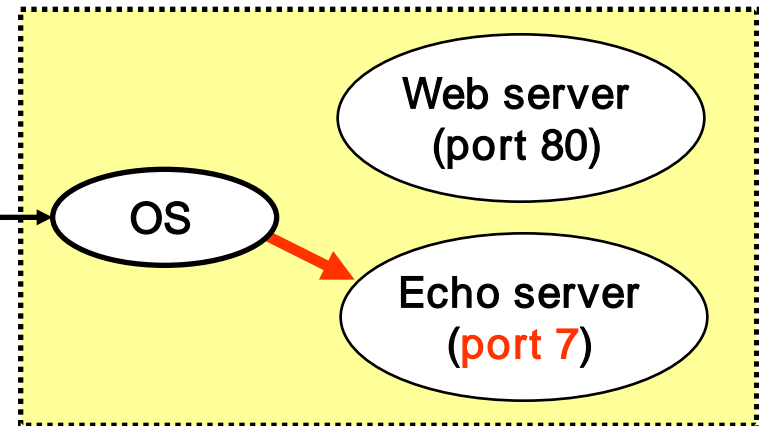
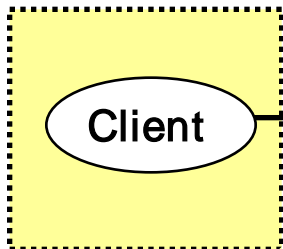
Server host **128.2.194.242**

Client host

Service request for
128.2.194.242:80
(i.e., the Web server)



Service request for
128.2.194.242:7
(i.e., the echo server)



Knowing What Port Number To Use



- Popular applications have well-known ports
 - E.g., port 80 for Web and port 25 for e-mail
 - See <http://www.iana.org/assignments/port-numbers>
- Well-known vs. ephemeral ports
 - Server has a well-known port (e.g., port 80)
 - Between 0 and 1023
 - Client picks an unused ephemeral (i.e., temporary) port
 - Between 1024 and 65535
- Uniquely identifying the traffic between the hosts
 - Two IP addresses and two port numbers
 - Underlying transport protocol (e.g., TCP or UDP)

Port Numbers are Unique on Each Host



- Port number uniquely identifies the socket
 - Cannot use same port number twice with same address
 - Otherwise, the OS can't demultiplex packets correctly
- Operating system enforces uniqueness
 - OS keeps track of which port numbers are in use
 - Doesn't let the second program use the port number
- Example: two Web servers running on a machine
 - They cannot both use port "80", the standard port #
 - So, the second one might use a non-standard port #
 - E.g., <http://www.cnn.com:8080>



UNIX Socket API

UNIX Socket API



- **Socket interface**
 - Originally provided in Berkeley UNIX
 - Later adopted by all popular operating systems
 - Simplifies porting applications to different OSes (even to the Windows!)
- **In UNIX, everything is like a file**
 - All input is like reading a file
 - All output is like writing a file
 - File is represented by an integer file descriptor
- **API implemented as system calls**
 - E.g., connect, read, write, close, ...

Typical Client Program



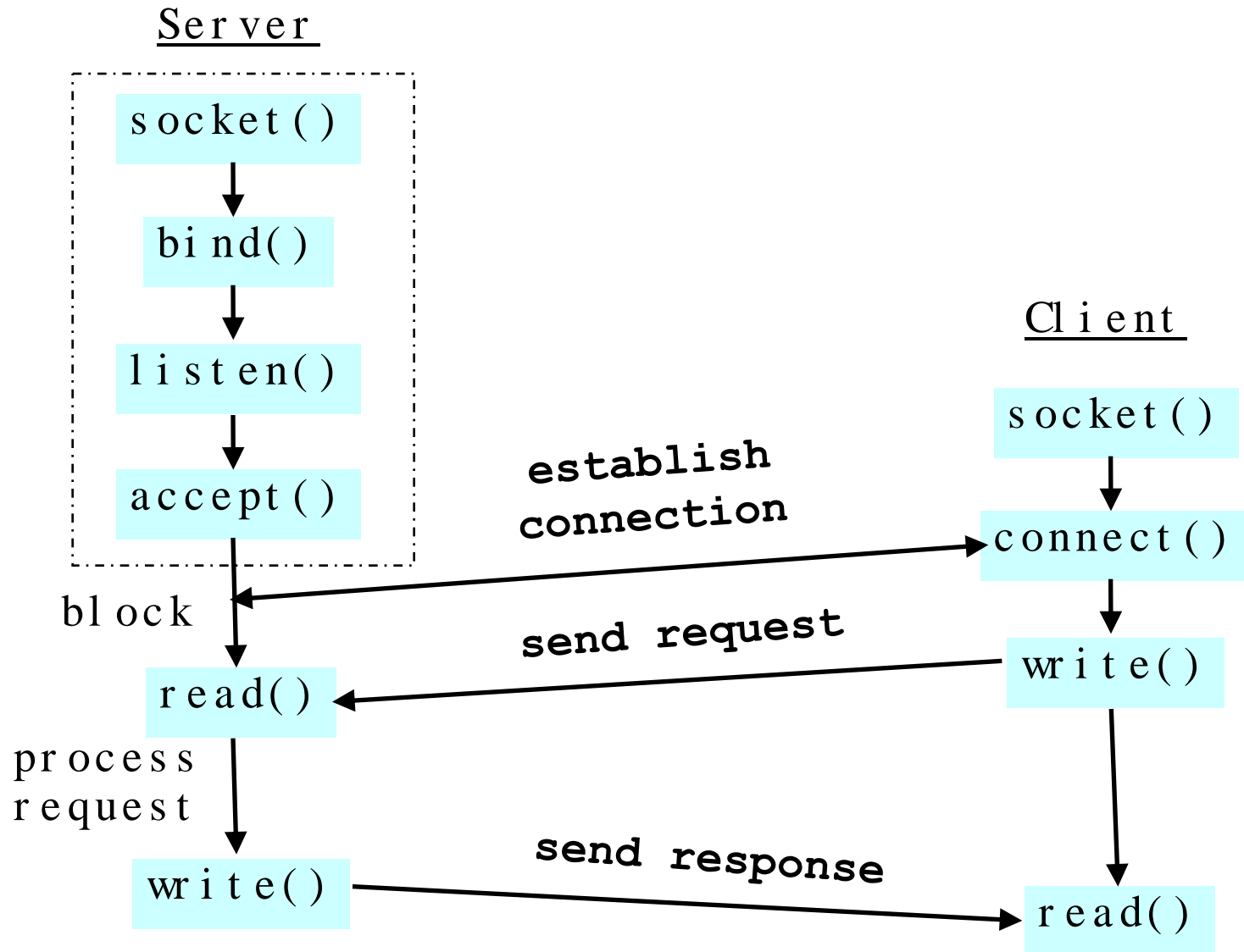
- Prepare to communicate
 - Create a socket
 - Determine server address and port number
 - Initiate the connection to the server
- Exchange data with the server
 - Write data to the socket
 - Read data from the socket
 - Do stuff with the data (e.g., render a Web page)
- Close the socket

Typical Server Program



- Prepare to communicate
 - Create a socket
 - Associate local address and port with the socket
- Wait to hear from a client (passive open)
 - Indicate how many clients-in-waiting to permit
 - Accept an incoming connection from a client
- Exchange data with the client over new socket
 - Receive data from the socket
 - Do stuff to handle the request (e.g., get a file)
 - Send data to the socket
 - Close the socket
- Repeat with the next connection request

Putting it All Together



Wanna See Real Clients and Servers?



- Apache Web server
 - Open source server first released in 1995
 - Name derives from “a patchy server” ;-)
 - Software available online at <http://www.apache.org>
- Mozilla Web browser
 - <http://www.mozilla.org/developer/>
- Sendmail
 - <http://www.sendmail.org/>
- BIND Domain Name System (Datagram)
 - Client resolver and DNS server
 - <http://www.isc.org/index.pl?/sw/bind/>
- ...



Wanna to have fun? Okay...

Client Programming

Client Creating a Socket: `socket()`



```
int socket(int domain, int type, int protocol)
```

- **Operation to create a socket**
 - Returns a descriptor (or handle) for the socket
 - Originally designed to support any protocol suite
- **Domain: protocol family**
 - PF_INET for the Internet
- **Type: semantics of the communication**
 - SOCK_STREAM: reliable byte stream
 - SOCK_DGRAM: message-oriented service
- **Protocol: specific protocol**
 - UNSPEC: unspecified
 - (PF_INET and SOCK_STREAM already implies TCP)

Client: Learning Server Address/Port



- Server typically known by name and service
 - E.g., “www.cnn.com” and “http”
- Need to translate into IP address and port #
 - E.g., “64.236.16.20” and “80”
- Translating the server’s name to an address
struct hostent *gethostbyname(char *name)
 - Argument: host name (e.g., “www.cnn.com”)
 - Returns a structure that includes the host address
- Identifying the service’s port number
struct servent *getservbyname(char *name, char *proto)
 - Arguments: service (e.g., “ftp”) and protocol (e.g., “tcp”)

Address/Port Data Structures



- Data Type: **struct hostent** - represent an entry in hosts database. Containing:
 - char *h_name
 - This is the “official” name of the host
 - char **h_aliases
 - These are alternative names for the host, represented as a null-terminated vector of strings
 - int h_addrtype
 - This is the host address type; in practice, its value is always either AF_INET or AF_INET6
 - int h_length
 - This is the length, in bytes, of each address
 - char **h_addr_list
 - This is the vector of addresses for the host
 - char *h_addr
 - This is a synonym for h_addr_list[0]

Address/Port Data Structures



- Data Type: **struct servent** - holds information about entries from services database. Containing:
 - char *s_name
 - This is the "official" name of the service
 - char **s_aliases
 - These are alternate names for the service, represented as an array of strings
 - int s_port
 - This is the port number for the service. Port numbers are given in network byte order
 - char *s_proto
 - This is the name of the protocol to use with this service

Client: Connecting Socket to the Server



```
int connect(int sockfd, struct sockaddr *server_address,  
            socklen_t addrlen)
```

- Client contacts the server to establish connection
 - Associate the socket with the server address/port
 - Acquire a local port number (assigned by the OS)
 - Request connection to server, who will hopefully accept
- Establishing the connection
 - Arguments: socket descriptor, server address, and address size
 - Returns 0 on success, and -1 if an error occurs

IP Address Data Structures



```
include <netinet/in.h>

// All pointers to socket address structures are often cast to pointers
// to this type before use in various functions and system calls:

struct sockaddr {
    unsigned short    sa_family;    // address family, AF_xxx
    char              sa_data[14]; // 14 bytes of protocol address
};

// IPv4 AF_INET sockets:

struct sockaddr_in {
    short            sin_family;    // e.g. AF_INET, AF_INET6
    unsigned short   sin_port;     // e.g. htons(3490)
    struct in_addr   sin_addr;     // see struct in_addr, below
    char             sin_zero[8];  // zero this if you want to
};

struct in_addr {
    unsigned long    s_addr;       // load with inet_pton()
};
```

Client: Sending and Receiving Data



- **Sending data**

`ssize_t write(int sockfd, void *buf, size_t len)`

- Arguments: socket descriptor, pointer to buffer of data to send, and length of the buffer
- Returns the number of characters written, and -1 on error

- **Receiving data**

`ssize_t read(int sockfd, void *buf, size_t len)`

- Arguments: socket descriptor, pointer to buffer to place the data, size of the buffer
- Returns the number of characters read (where 0 implies “end of file”), and -1 on error

- **Closing the socket**

`int close(int sockfd)`



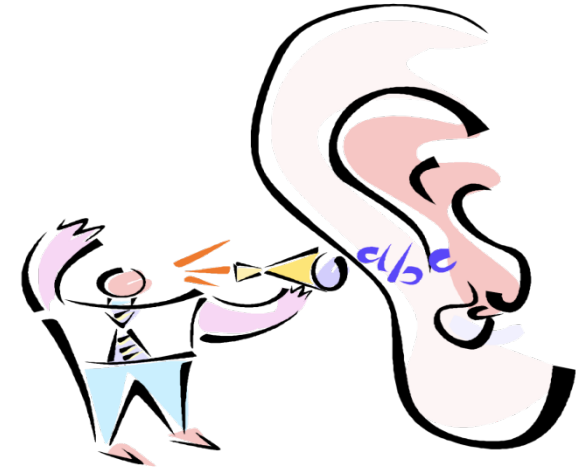
Not enough fun? Okay... face a headache!

Server Programming

Servers Differ From Clients



- **Passive open**
 - Prepare to accept connections
 - ... but don't actually establish
 - ... until hearing from a client
- **Hearing from multiple clients**
 - Allowing a backlog of waiting clients
 - ... in case several try to communicate at once
- **Create a socket for each client**
 - Upon accepting a new client
 - ... create a new socket for the communication



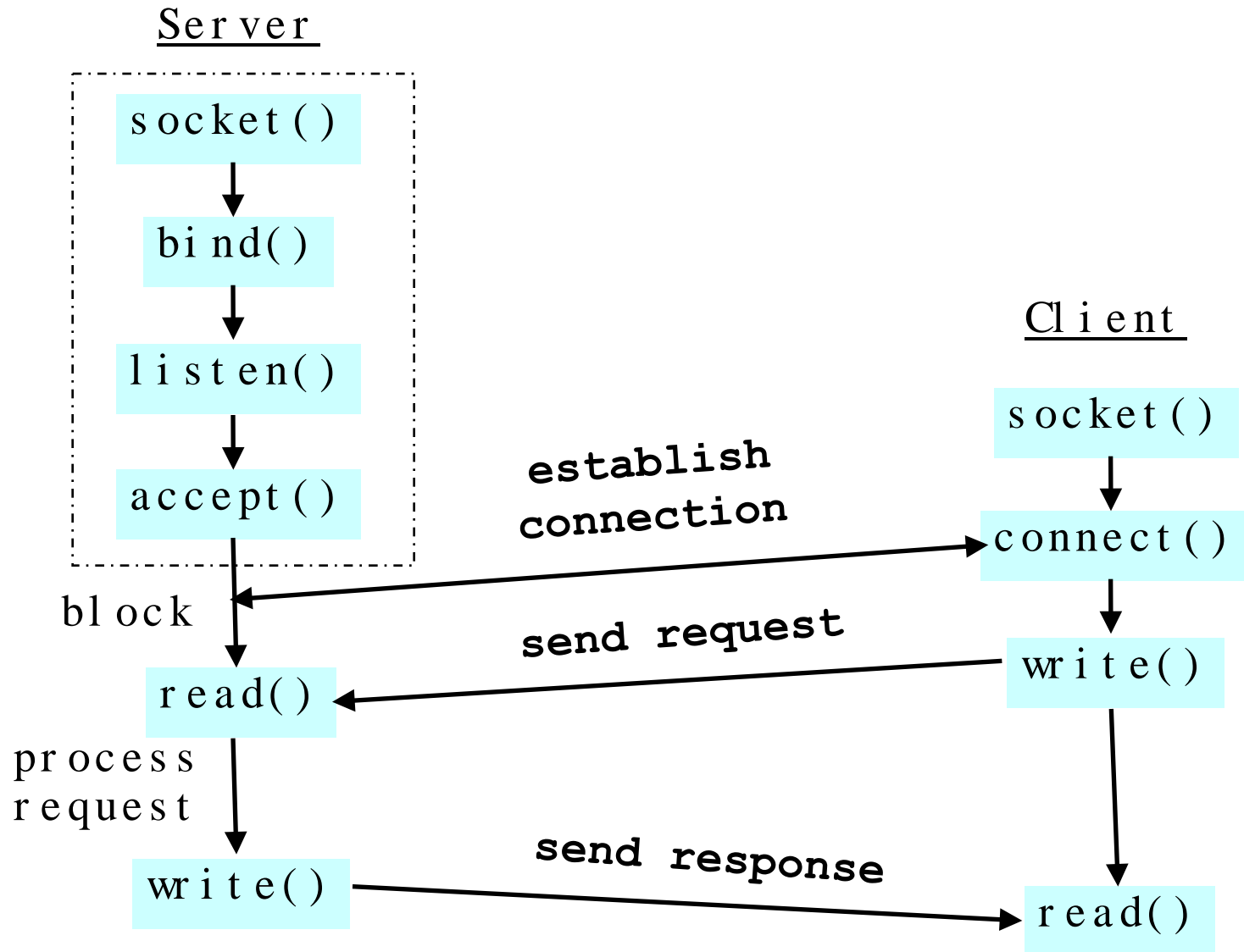
Remember: Typical Server Program



- Prepare to communicate
 - Create a socket
 - Associate local address and port with the socket
- Wait to hear from a client (passive open)
 - Indicate how many clients-in-waiting to permit
 - Accept an incoming connection from a client
- Exchange data with the client over new socket
 - Receive data from the socket
 - Do stuff to handle the request (e.g., get a file)
 - Send data to the socket
 - Close the socket
- Repeat with the next connection request



Remember: The Big Picture



Server: Server Preparing its Socket



- **Server creates a socket and binds address/port**
 - Server creates a socket, just like the client does
 - Server associates the socket with the port number (and hopefully no other process is already using it!)
- **Create a socket**
`int socket(int domain, int type, int protocol)`
- **Bind socket to the local address and port number**
`int bind(int sockfd, struct sockaddr *my_addr, socklen_t addrlen)`
 - Arguments: socket descriptor, server address, address length
 - Returns 0 on success, and -1 if an error occurs

Server: Allowing Clients to Wait



- Many client requests may arrive
 - Server cannot handle them all at the same time
 - Server could reject the requests, or let them wait
 - Define how many connections can be pending: backlog
- Wait for clients

```
int listen(int sockfd, int backlog)
```

 - Arguments: socket descriptor and acceptable backlog
 - Returns a 0 on success, and -1 on error
- What if too many clients arrive?
 - Some requests don't get through
 - The Internet makes no promises...
 - And the client can always try again



Server: Accepting Client Connection



- Now all the server can do is wait...
 - Waits for connection request to arrive
 - Blocking until the request arrives
 - And then accepting the new request



- Accept a new connection from a client

```
int accept(int sockfd, struct sockaddr *addr, socketlen_t  
*addrlen)
```

- Arguments: socket descriptor, structure that will provide client address and port, and length of the structure
- Returns descriptor for a new socket for this connection

Server: One Request at a Time?



- Serializing requests is inefficient
 - Server can process just one request at a time
 - All other clients must wait until previous one is done
- May need to time share the server machine
 - Alternate between servicing different requests
 - E.g. use multi-threading
 - Or, start a new process to handle each request
 - Allow the operating system to share the CPU across processes
 - Or, some hybrid of these two approaches

Client and Server: Cleaning House



- Once the connection is open
 - Both sides can read and write
 - Two unidirectional streams of data
 - In practice, client writes first, and server reads
 - ... then server writes, and client reads, and so on
- Closing down the connection
 - Either side can close the connection
 - ... using the `close()` system call
- What about the data still “in flight”
 - Data in flight still reaches the other end
 - So, server can `close()` before client finishing reading



The Problem of Interoperability

Byte Order



- Hosts differ in how they store data
 - E.g., four-byte number (byte3, byte2, byte1, byte0)
- Little endian (“little end comes first”) ← Intel PCs!!!
 - Low-order byte stored at the lowest memory location
 - Byte0, byte1, byte2, byte3
- Big endian (“big end comes first”)
 - High-order byte stored at lowest memory location
 - Byte3, byte2, byte1, byte 0
- Makes it more difficult to write portable code
 - Client may be big or little endian machine
 - Server may be big or little endian machine

IP is Big Endian



- But, what byte order is used “on the wire”
 - That is, what do the network protocols use?
- The Internet Protocols picked one convention
 - IP is big endian (aka “network byte order”)
- Writing portable code requires conversion
 - Use `htons()` and `htonl()` to convert to network byte order
 - Use `ntohs()` and `ntohl()` to convert to host order
- Hides details of what kind of machine you’re on
 - Use the system calls when sending/receiving data structures longer than one byte

Why Can't Sockets Hide These Details?



- Dealing with endian differences is tedious
 - Couldn't the socket implementation deal with this
 - ... by swapping the bytes as needed?
- No, swapping depends on the data type
 - Two-byte short int: (byte 1, byte 0) vs. (byte 0, byte 1)
 - Four-byte long int: (byte 3, byte 2, byte 1, byte 0) vs. (byte 0, byte 1, byte 2, byte 3)
 - String of one-byte characters: (char 0, char 1, char 2, ...) in both cases
- Socket layer doesn't know the data types
 - Sees the data as simply a buffer pointer and a length
 - Doesn't have enough information to do the swapping



The Web as an Example Application

The Web: URL, HTML, and HTTP



- **Uniform Resource Locator (URL)**
 - A pointer to a “black box” that accepts request methods
 - Formatted string with protocol (e.g., http), server name (e.g., www.cnn.com), and resource name (coolpic.jpg)
- **HyperText Markup Language (HTML)**
 - Representation of hypertext documents in ASCII format
 - Format text, reference images, embed hyperlinks
 - Interpreted by Web browsers when rendering a page
- **HyperText Transfer Protocol (HTTP)**
 - Client-server protocol for transferring resources
 - Client sends request and server sends response



URI vs. URL

- Uniform Resource Identifier (URI)
 - Addresses a resource including the
 - Protocol
 - Machine address
 - The path
- Uniform Resource Locator (URL)
 - An special URI
 - Addresses a resource via a representation of its primary access mechanism
 - Each URL is URI too, but not vice versa
 - Each resource could be identified via
 - Many URI addresses
 - But only one URL address (usually the network address)

Example: HyperText Transfer Protocol



```
GET /courses/archive/spring08/cos461/ HTTP/1.1
```

```
Host: www.cs.princeton.edu
```

```
User-Agent: Mozilla/4.03
```

```
<CRLF>
```

Request

```
HTTP/1.1 200 OK
```

```
Date: Mon, 4 Feb 2008 13:09:03 GMT
```

```
Server: Netscape-Enterprise/3.5.1
```

```
Content-Type: text/plain
```

```
Last-Modified: Mon, 4 Feb 2008 11:12:23 GMT
```

```
Content-Length: 21
```

```
<CRLF>
```

```
Site under construction
```

Response

Example Client: Web Browser



- **Generating HTTP requests**
 - User types URL, clicks a hyperlink, or selects bookmark
 - User clicks “reload”, or “submit” on a Web page
 - Automatic downloading of embedded images
- **Layout of response**
 - Parsing HTML and rendering the Web page
 - Invoking helper applications (e.g., Acrobat, PowerPoint)
- **Maintaining a cache**
 - Storing recently-viewed objects
 - Checking that cached objects are fresh

Client: Typical Web Transaction



- User clicks on a hyperlink
 - `http://www.cnn.com/index.html`
- Browser learns the IP address
 - Invokes `gethostbyname(www.cnn.com)`
 - And gets a return value of `64.236.16.20`
- Browser creates socket and connects to server
 - OS selects an ephemeral port for client side
 - Contacts `64.236.16.20` on port 80
- Browser writes the HTTP request into the socket
 - “`GET /index.html HTTP/1.1`
`Host: www.cnn.com`
`<CRLF>`”

In Fact, Try This at a UNIX Prompt...



```
l abpc: telnet www.cnn.com 80
GET /index.html HTTP/1.1
Host: www.cnn.com
<CRLF>
```

And you'll see the response...

Client: Typical Web Transaction (Cont)



- Browser parses the HTTP response message
 - Extract the URL for each embedded image
 - Create new sockets and send new requests
 - Render the Web page, including the images
- Opportunities for caching in the browser
 - HTML file
 - Each embedded image
 - IP address of the Web site

Web Server



- Web site vs. Web server
 - **Web site**: collections of Web pages associated with a particular host name
 - **Web server**: program that satisfies client requests for Web resources
- Handling a client request
 - Accept the socket
 - Read and parse the HTTP request message
 - Translate the URL to a filename
 - Determine whether the request is authorized
 - Generate and transmit the response

Conclusions



- **Client-server paradigm**
 - Model of communication between end hosts
 - Client asks, and server answers
- **Sockets**
 - Simple byte-stream and messages abstractions
 - Common application programmable interface
- **HyperText Transfer Protocol (HTTP)**
 - Client-server protocol
 - URL, HTML, and HTTP
- **A Good Online Tutorial**
 - Beej's Guide to Network Programming